JP,2000-261499,A

JAPANESE [JP,2000-261499,A]

CLAIMS DETAILED DESCRIPTION TECHNICAL FIELD EFFECT OF THE INVENTION MEANS DESCRIPTION OF DRAWINGS DRAWINGS



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CLAIMS

[Claim(s)]

[Claim 1] It is the approach performed in the equipment which performs data transmission and reception between external terminals, using either of two or more channels alternatively. Between a part or all of two or more of said channels, and said equipment Make the proxy equipment which executes processing about said data transmission and reception by proxy, respectively intervene, and a logical mirror site is built. By performing a communication link local between each proxy equipment and said equipment, and communicating the gestalt depending on the channel connected to the proxy equipment concerned between each proxy equipment and said external terminal The traffic dispersion approach characterized by distributing the traffic in the equipment concerned.

[Claim 2] It is the approach performed in the equipment which transmits response data to the addressing to an external terminal concerned through either of two or more channels when a data Request to Send is received from an external terminal. Relate each of two or more of said channels with the predetermined data characteristic, and the data characteristic of the response data concerned is specified in the case of transmission of the response data to said external terminal. The traffic dispersion approach characterized by distributing the traffic in said equipment by leading said response data to the channel which suits this specified data characteristic.

[Claim 3] the traffic dispersion approach according to claim 2 characterize by said predetermined data characteristic be the information determine according to each engine performance or quality of a channel including at least one of the size of the communications protocol use between said external terminals, and said response data, the contents of data of a specific location, and the priority **s of a response.

[Claim 4] The traffic dispersion approach according to claim 2 or 3 characterized by specifying the data characteristic of said response data which should be transmitted based on the contents of said data Request to Send which received.

[Claim 5] The traffic dispersion approach according to claim 2 or 3 which the proxy equipment which executes processing about said data transmission and reception by proxy, respectively is made to intervene between a part or all of two or more of said channels, and said equipment, and is characterized by performing transmission of the response data to said external terminal through said proxy equipment at least.

[Claim 6] In the communication system equipped with the communication device which performs data transmission and reception between external terminals, using either of two or more channels alternatively Between a part or all of two or more of said channels, and said communication device, arrange the proxy equipment which executes processing about said data transmission and reception by proxy, respectively, and a logical mirror site is built with each proxy equipment and said communication device to it. Communication system which performs a communication link local between each proxy equipment and said communication device, and is characterized by constituting between each proxy equipment and said external terminal so that the gestalt depending on the channel connected to the proxy equipment concerned may be communicated.

[Claim 7] In the communication system which transmits response data to the addressing to an external terminal concerned through either of two or more channels when a data Request to Send is received from an external terminal Two or more interface devices for corresponding to each of two or more of said channels, and 1 to 1, and enabling the communication link between each channels, While holding the correlation information on a distinction means to distinguish the data characteristic of response data based on said data Request to Send, and each of two or more of said interface devices and a predetermined data characteristic Communication system characterized by forming the management tool which specifies the interface device corresponding to the data characteristic distinguished with said distinction means based on said correlation information, and leads said response data to said specified interface device. [Claim 8] It is the communication system according to claim 7 characterized by said management tool specifying the interface device connected to the channel based on said distinguished communications protocol by said distinction means distinguishing the communications protocol used between said external terminals as said data characteristic. [Claim 9] It is the communication system according to claim 7 with which it is characterized by said management tool specifying the interface device by which said distinguished size is connected to descending at a mass channel by said distinction means distinguishing the size of the size of said response data as said data characteristic.

[Claim 10] It is the communication system according to claim 7 characterized by said management tool specifying the interface device connected to the channel beforehand assigned for said every contents of the distinguished data by said distinction means distinguishing the contents of data of the specific location of said response data as said data characteristic. [Claim 11] It is the communication system according to claim 7 characterized by said management tool specifying the interface device connected to a mass channel at order with said distinguished high priority by said distinction means distinguishing the priority of the response corresponding to said data Request to Send as said data characteristic.

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DETAILED DESCRIPTION

[Detailed Description of the Invention]

[The technical field to which invention belongs] This invention relates to the technique which controls concentration of the traffic in the server connected to the Internet. [0002]

[Background of the Invention] It is common that service which performs information offer by WWW (World Wide Web) of the Internet, FTP (File Transfer Protocol), etc. is offered through a network service provider (it is called a "provider" for short Network Service Provider and the following). Service which a provider offers is offered when a user mainly accesses at arbitration the server which accumulates information. For this reason, if access from a user concentrates on a provider's server, in the channel which that provider is preparing, delay of traffic may occur or delay of traffic may occur in the channel between the servers and providers who are accumulating information. When such delay occurs, enjoyment of smooth information becomes impossible to a user.

[0003] Then, in order to cancel delay of traffic, it is performed that both users make two or more servers distribute traffic by accumulating the same accessible information ("mirror information") conventionally. Generally the location which offers mirror information is called the "mirror site." The mirror site serves as two or more sites and a gestalt to which two sites 100 and 110 were connected more simply at the Internet 10, as shown in drawing 6. In the server 101 and the site 110, the server 111 is arranged by the site 100, respectively. Mirror information is accumulated in each servers 101 and 111, respectively, and the interface devices 102 and 112 for controlling connection between an internal network and channels 20 and 21 etc. are arranged further. By preparing such a mirror site, it is lost that access from a user concentrates on one server, and distribution of traffic is attained. However, if a mirror site is prepared, the cost which at least two lines [at least two], i.e., a server, are needed, and construction of a system and a maintenance take to the independent site (server) will become large.

[0004] The gestalt which distributes the traffic in the server concerned can be considered by preparing the interface device which can connect either of two or more channels to one set of a server alternatively as a solution means of such a problem, for example, a router. That is, as shown in drawing 7, the interface devices 211 and 212 for connecting self and the Internet 10 to one set of the server 210 currently prepared for the site 200 through either of two channels 20 and 21 are arranged, and traffic is distributed.

[0005] However, traffic is not necessarily well distributed by channels 20 and 21 only by arranging two or more interface devices 211 and 212 like illustration. In case a server 210 transmits data according to a demand of a user, this is because the suitable interface device cannot be specified surely, in order not to memorize the information on the passed data, when many of operating systems (Operating System) used for a server 210 receive the data transmitted by the user.

[0006] That is, in a server equipped with a router, it is decided whether to transmit data using which channel by the contents of record of the routing table with which a server is equipped. For example, if it is data transmission to the communication device with which IP (internet protocol)

address which a server 210 manages was assigned, with reference to routing table, the channel connected to the communication device can be specified. However, when transmitting the data which answer a data Request to Send from the communication device with which the unspecified address, for example, a global address, was assigned, since the IP address of the communication device is unrecordable on routing table, a site 210 cannot specify the channel (that is, interface device) used in case the data which answer are transmitted. For this reason, traffic cannot be distributed well.

[0007] Then, this invention makes it a technical problem to offer the traffic control approach which enables distribution of traffic by simple technique. Other technical problems of this invention are to offer the cheap communication system equipped with the traffic dispersion function.

[8000]

[Means for Solving the Problem] The traffic control approach of the 1st invention which solves the above-mentioned technical problem It is the approach performed in the equipment which performs data transmission and reception between external terminals, using either of two or more channels alternatively. Between a part or all of two or more of said channels, and said equipment Make the proxy equipment which executes processing about said data transmission and reception by proxy, respectively intervene, and a logical mirror site is built. Between each proxy equipment and said equipment, it is characterized by distributing the traffic in the equipment concerned by performing a local communication link and communicating the gestalt depending on the channel connected to the proxy equipment concerned between each proxy equipment and said external terminal.

[0009] The traffic control approach of the 2nd invention is an approach performed in the equipment which transmits response data to the addressing to an external terminal concerned through either of two or more channels when a data Request to Send is received from an external terminal. Relate each of two or more of said channels with the predetermined data characteristic, and the data characteristic of the response data concerned is specified in the case of transmission of the response data to said external terminal. By leading said response data to the channel which suits this specified data characteristic, it is characterized by distributing the traffic in said equipment.

[0010] Here, a predetermined data characteristic is information determined according to each engine performance or quality of a channel including at least one of the size of the communications protocol used between for example, the aforementioned external terminals, and said response data, the contents of data of a specific location, and the priority **s of a response. This data characteristic may be specified from the header of response data etc., or you may make it specify it based on the contents of said data Request to Send which received. Moreover, the proxy equipment which executes processing about said data transmission and reception by proxy, respectively is made to intervene between a part or all of two or more of said channels, and said equipment, and it may be made to perform transmission of the response data to said external terminal through said proxy equipment at least.

[0011] The communication system of this invention which solves a technical problem besides the above In the communication system equipped with the communication device which performs data transmission and reception between external terminals, using either of two or more channels alternatively Arrange the proxy equipment which executes processing about said data transmission and reception by proxy, respectively between a part or all of two or more of said channels, and said communication device, and a logical mirror site is built with each proxy equipment and said communication device to it. Between each proxy equipment and said communication link is performed, and between each proxy equipment and said external terminal, it constitutes so that the gestalt depending on the channel connected to the proxy equipment concerned may be communicated. Said proxy equipment is changing the address of a data transmission place into the transmission place address of self while changing the address of for example, data transmitting origin into the transmitting agency address of self, and it forms the data transceiver gestalt between said external terminals and said communication devices virtually.

[0012] In the communication system which transmits response data to the addressing to an external terminal concerned through either of two or more channels when the communication system concerning other configurations of this invention receives a data Request to Send from an external terminal Two or more interface devices for corresponding to each of two or more of said channels, and 1 to 1, and enabling the communication link between each channels, While holding the correlation information on a distinction means to distinguish the data characteristic of response data based on said data Request to Send, and each of two or more of said interface devices and a predetermined data characteristic The interface device corresponding to the data characteristic distinguished with said distinction means is specified based on said correlation information, and it is characterized by forming the management tool which leads said response data to said specified interface device.

[0013]

[Embodiment of the Invention] Hereafter, the gestalt of operation of this invention is explained with reference to a drawing.

(The 1st operation gestalt) Drawing 1 is the important section block diagram of the communication system which applied this invention. Among the external terminals (a "client" is called hereafter) 30 and 31 connected to the wide area network 10, for example, the Internet, respectively, a site 300 connects mutually the communication device (a "server" is called hereafter) 310 in which two-way communication is possible, and the proxy server 320 used as an example of proxy (proxy) equipment through an internal network, and is constituted. [0014] The server 310 is equipped with the interface device 311 to which the private address b was assigned, and the interface device 312 to which the global address F was assigned. [0015] It is the substitute communication device equipped with the function to execute transmission and reception of the data between clients 30 and 31 and a server 310 by proxy, a proxy server 320 is changing the address of a data transmission place into the transmission place address of self while it changes the address of data transmitting origin into the transmitting agency address of self, and it forms the data transceiver gestalt between clients 30 and 31 and a server 310 virtually. This proxy server 320 is equipped with the interface device 321 to which the private address c was assigned, and the interface device 322 to which the global address D was assigned.

[0016] In the example of <u>drawing 1</u>, interface devices 311 and 321 perform control to control the internal two-way communication between the server 310-proxy servers 320 through an internal network (local communication link), and for interface devices 312 and 322 each, connect a server 310 and a proxy server 320 to the Internet 10, respectively. In addition, in the communication system of this operation gestalt, the data transmitted and received between a server 310 and each clients 30 and 31 shall be divided under a fixed regulation, and shall be performed in the form of [which the destination label attached] a packet. Hereafter, the data transmitted and received may be called a "packet."

[0017] Actuation of the communication system of this operation gestalt is as follows. Here, the case where the client 30 to which Address A was assigned, and the client 31 to which Address E was assigned publish a data Request to Send towards a server 310 at a coincidence term is explained.

[0018] The data Request to Send from a client 30 shall be published to the proxy server 320 which has an interface device 322. A proxy server 320 will publish a data Request to Send to a server 310 instead of a client 30, if this data Request to Send is received. The publishing agency address of the data Request to Send at this time turns into Address c. That is, the address A of the client 30 which is the issue origin of a data Request to Send from the first is changed into the address c of a proxy server 320, and is sent to a server 310. Although a server 310 transmits a response packet corresponding to this data Request to Send, the transmission place of the response packet at this time serves as a proxy server 320. Once a proxy server 320 changes into Address c the address b by the side of the server 310 which transmitted the response packet, it changes this into the address D of an interface device 322, and transmits the response packet which includes this address D in a header unit to a client 30.

[0019] A server 310 does not need to be conscious of the address A of the issue origin (client

30) of a data Request to Send, and just comes to answer the proxy server 320 by adopting such a gestalt in the data (packet) demanded from the proxy server 320. On the other hand, direct access of the client 31 can be carried out to a server 310 through the interface device 312 with a channel 21 and Address F.

[0020] That is, from the Internet 10 side (a client 30 side, 31 sides), it is visible as if the server with Address D and the server with Address F existed independently, respectively. On the other hand, a local communication link is performed between a server 310 and a proxy server 320. Consequently, the mirror site which has two servers virtually can be built by one set of a server 310, and it becomes possible to distribute the traffic in a server 310 simply and suitable for two channels 20 and 21. Moreover, since it is not necessary to install two sets of servers, a manufacturing cost and maintenance employment expense become cheap.

[0021] (The 2nd operation gestalt) <u>Drawing 2</u> is drawing having shown the example of a configuration of the communication system which formed the above-mentioned proxy server 320 and the proxy servers 420 and 430 of the same function between a server 410 and channels 20 and 21, respectively. One site where the sign 400 arranged this communication system, the interface device (the interface devices 311 and 312 and this function of the 1st operation gestalt) which a server 410 equips with 411 and 412, and 421, 422, 431 and 432 are interface devices (the interface devices 321 and 322 and this function of the 1st operation gestalt) with which each proxy servers 420 and 430 are equipped.

[0022] With the communication system of this operation gestalt as well as the communication system of the 1st operation gestalt, two servers can be virtually built by the single server 410. The other party of data transmission and reception of the server 410 in this case always becomes proxy servers 420 and 430. Consequently, it becomes possible to distribute simply and appropriately the traffic in the server 410 when there is a data Request to Send from two clients 30 and 31.

[0023] (The 3rd operation gestalt) Drawing 3 is the block diagram of the communication system which can distribute traffic, without making a proxy server intervene. Although the channels 20 and 21 which constitute a part use the Internet 10 the same as the thing of the 1st and 2nd operation gestalt and it is not illustrated for convenience, two or more clients shall be connected to the Internet 10. The server 510 used as an example of the communication system in this case equips each of two channels 20 and 21 with the interface devices 511 and 512 corresponding to 1 to 1. The respectively global address is assigned to these interface devices 511 and 512. [0024] First, the structure of the packet (a data Request to Send, response packet) used in this operation gestalt is explained. Data characteristics, such as a transmitting agency, and a communications protocol besides the address of a transmission place, are usually included in the header unit of this kind of packet as additional information. Drawing 4 (a) - (c) is drawing which illustrated the example of a frame structure of a response packet, and the response packet to which Packets A and C are transmitted through an interface device 511 (address 255.20.34.12), and Packet B are packets transmitted through an interface device 512 (address: 255.20.34.13). The transmission place of these packet A-C serves as a client (address: 155.19.25.37, 175.20.35.10, 185.45.67.21) different, respectively, and, as for HTTP and Packet B, also in the classification of a communications protocol, Packets A and C serve as FTP. [0025] The thing [priority] according to a packet is recorded. A "priority" is an index which shows the significance (or quality) of the response for example, according to the contents of service. That is, it is not to concern a server 510 with reception sequence, but transmit a response packet sequentially from what has a high priority to the data Request to Send received within a fixed period. In addition, the same header unit is added also about the data Request to Send, and it is used in the various distinction sections 513-516 mentioned later for the specification of a data characteristic.

[0026] The output interface management section 17 performs control for choosing either of the interface devices 511 and 512 at the time of management of routing table 518, and transmission of the response packet according to the data Request to Send from a client, and drawing a response packet at it. Selection of interface devices 511 and 512 is performed based on the newest contents of record of routing table 518 in principle.

[0027] Routing table 518 records the correlation information on each of interface devices 511 and 512, and a predetermined data characteristic in consideration of the engine performance of channels 20 and 21, quality, etc. <u>Drawing 5</u> is drawing having shown the example of the contents of record of this routing table 518, and the data characteristic distinguished in the various distinction sections 513–516 mentioned later besides the transmission place address and the transmitting agency address is recorded. An "output interface device" is an interface device beforehand associated corresponding to each data characteristic.

[0028] The protocol distinction section 513 distinguishes the communications protocol used between clients as a data characteristic. If this communications protocol is FTP and HTTP, it is a mass channel and a data characteristic to choose the channel of small capacity, if it becomes other than this. The packet size distinction section 514 distinguishes the size of a packet size (data size) as a data characteristic. A packet size is a data characteristic to choose a mass channel, so that it is long. The output priority distinction section 515 distinguishes the priority of the response for example, corresponding to a data Request to Send as a data characteristic. This priority determines the sequence of a response according to the importance (or height of quality) of service as mentioned above, and also the thing which has a high priority can use it as a data characteristic to choose a more nearly mass channel. 1 byte of low order of a destination address [in / in the data bit distinction section 516 / the contents of data of the specific location of data, for example, a header unit,] distinguishes odd number or even number as said data characteristic.

[0029] These data characteristics can usually be distinguished in analyzing the contents of the header unit in the case of a data Request to Send. This is because the response packet of the contents according to the data Request to Send will be generated. However, you may make it distinguish the above-mentioned data characteristic from the header unit of a response packet. The distinction result by each distinction sections 513-516 is notified to the output interface management section 517, respectively, and the contents are recorded on routing table 518. In addition, the contents of the above-mentioned data characteristic are instantiation, and are not limited to these things.

[0030] Thus, in the communication system constituted, actuation when there is a data Request to Send from a certain client is as follows. The information on the header unit of a data Request to Send is sent to each distinction sections 513–516, and the data characteristic of the response packet which corresponds, respectively is distinguished. A distinction result is recorded on routing table 518 through the output interface management section 517. If a response packet is generated by the response packet generation means which is not illustrated, the output interface management section 517 will determine to which of interface devices 511 and 512 this response packet is led.

[0031] Although the contents of record of routing table 518 are followed as mentioned above in principle on the occasion of decision, it can also be based only on the distinction result of each distinction sections 513–516. For example, if the number of the tails of the address of a transmission place is odd, in the case of an interface device 511 and even number, an interface device 512 will be determined as an interface device for transmission of a response packet. Or if the protocol of a response packet is HTTP or FTP, an interface device 511 and when other, an interface device 512 will be determined as an interface device for transmission of a response packet. Or when a priority is "quantity" and an interface device 511 and a transmitting priority are "low", an interface device 512 is determined as an interface device for transmission of a response packet. Or the method of determining an interface device by combining suitably the address of the decision and the transmission place according to the merits and demerits of response packet length, a communications protocol, a priority, data size, etc. is also one of the gestalten of operation. According to this communication system, the traffic in a server 510 can be appropriately distributed now.

[0032] the communication link which the response (turn around time) of a server will improve and will be performed through the Internet by building the communication system of the 1st – the 3rd operation gestalt as mentioned above if it sees from a user — breaking off (quality fluctuation) — the decreasing effectiveness is acquired, if it sees on the other hand from the

side which provides this communication system with service information, the installation number of a server becomes fewer and the effectiveness that the cost which plant—and—equipment investment takes can reduce remarkably will be acquired. In order to build a mirror site logically by one set of a server, it becomes unnecessary moreover, to take the synchronization of an updating stage when it becomes unnecessary to have reproduced the contents in the case of building a mirror site by two sets of servers and there is modification of contents.

[0033] In addition, although the 1st – the 3rd operation gestalt showed the example in case the number of channels is two, this invention is applicable similarly about the case (in this case, the interface device corresponding to 1 to 1 is needed for each channel) of three or more channels.

[Effect of the Invention] According to this invention, the traffic dispersion approach that distribution of traffic can be aimed at by simple technique can be offered now so that clearly from the above explanation. Moreover, the cheap communication system equipped with the traffic dispersion function can be offered.

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TECHNICAL FIELD

[The technical field to which invention belongs] This invention relates to the technique which controls concentration of the traffic in the server connected to the Internet. [0002]

[Background of the Invention] It is common that service which performs information offer by WWW (World Wide Web) of the Internet, FTP (File Transfer Protocol), etc. is offered through a network service provider (it is called a "provider" for short Network Service Provider and the following). Service which a provider offers is offered when a user mainly accesses at arbitration the server which accumulates information. For this reason, if access from a user concentrates on a provider's server, in the channel which that provider is preparing, delay of traffic may occur or delay of traffic may occur in the channel between the servers and providers who are accumulating information. When such delay occurs, enjoyment of smooth information becomes impossible to a user.

[0003] Then, in order to cancel delay of traffic, it is performed that both users make two or more servers distribute traffic by accumulating the same accessible information ("mirror information") conventionally. Generally the location which offers mirror information is called the "mirror site." The mirror site serves as two or more sites and a gestalt to which two sites 100 and 110 were connected more simply at the Internet 10, as shown in drawing 6. In the server 101 and the site 110, the server 111 is arranged by the site 100, respectively. Mirror information is accumulated in each servers 101 and 111, respectively, and the interface devices 102 and 112 for controlling connection between an internal network and channels 20 and 21 etc. are arranged further. By preparing such a mirror site, it is lost that access from a user concentrates on one server, and distribution of traffic is attained. However, if a mirror site is prepared, the cost which at least two lines [at least two], i.e., a server, are needed, and construction of a system and a maintenance take to the independent site (server) will become large.

[0004] The gestalt which distributes the traffic in the server concerned can be considered by preparing the interface device which can connect either of two or more channels to one set of a server alternatively as a solution means of such a problem, for example, a router. That is, as shown in drawing 7, the interface devices 211 and 212 for connecting self and the Internet 10 to one set of the server 210 currently prepared for the site 200 through either of two channels 20 and 21 are arranged, and traffic is distributed.

[0005] However, traffic is not necessarily well distributed by channels 20 and 21 only by arranging two or more interface devices 211 and 212 like illustration. In case a server 210 transmits data according to a demand of a user, this is because the suitable interface device cannot be specified surely, in order not to memorize the information on the passed data, when many of operating systems (Operating System) used for a server 210 receive the data transmitted by the user.

[0006] That is, in a server equipped with a router, it is decided whether to transmit data using which channel by the contents of record of the routing table with which a server is equipped. For example, if it is data transmission to the communication device with which IP (internet protocol) address which a server 210 manages was assigned, with reference to routing table, the channel connected to the communication device can be specified. However, when transmitting the data

which answer a data Request to Send from the communication device with which the unspecified address, for example, a global address, was assigned, since the IP address of the communication device is unrecordable on routing table, a site 210 cannot specify the channel (that is, interface device) used in case the data which answer are transmitted. For this reason, traffic cannot be distributed well.

[0007] Then, this invention makes it a technical problem to offer the traffic control approach which enables distribution of traffic by simple technique. Other technical problems of this invention are to offer the cheap communication system equipped with the traffic dispersion function.

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EFFECT OF THE INVENTION

[Effect of the Invention] According to this invention, the traffic dispersion approach that distribution of traffic can be aimed at by simple technique can be offered now so that clearly from the above explanation. Moreover, the cheap communication system equipped with the traffic dispersion function can be offered.

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MEANS

[Means for Solving the Problem] The traffic control approach of the 1st invention which solves the above-mentioned technical problem It is the approach performed in the equipment which performs data transmission and reception between external terminals, using either of two or more channels alternatively. Between a part or all of two or more of said channels, and said equipment Make the proxy equipment which executes processing about said data transmission and reception by proxy, respectively intervene, and a logical mirror site is built. Between each proxy equipment and said equipment, it is characterized by distributing the traffic in the equipment concerned by performing a local communication link and communicating the gestalt depending on the channel connected to the proxy equipment concerned between each proxy equipment and said external terminal.

[0009] The traffic control approach of the 2nd invention is an approach performed in the equipment which transmits response data to the addressing to an external terminal concerned through either of two or more channels when a data Request to Send is received from an external terminal. Relate each of two or more of said channels with the predetermined data characteristic, and the data characteristic of the response data concerned is specified in the case of transmission of the response data to said external terminal. By leading said response data to the channel which suits this specified data characteristic, it is characterized by distributing the traffic in said equipment.

[0010] Here, a predetermined data characteristic is information determined according to each engine performance or quality of a channel including at least one of the size of the communications protocol used between for example, the aforementioned external terminals, and said response data, the contents of data of a specific location, and the priority **s of a response. This data characteristic may be specified from the header of response data etc., or you may make it specify it based on the contents of said data Request to Send which received. Moreover, the proxy equipment which executes processing about said data transmission and reception by proxy, respectively is made to intervene between a part or all of two or more of said channels, and said equipment, and it may be made to perform transmission of the response data to said external terminal through said proxy equipment at least.

[0011] The communication system of this invention which solves a technical problem besides the above In the communication system equipped with the communication device which performs data transmission and reception between external terminals, using either of two or more channels alternatively Arrange the proxy equipment which executes processing about said data transmission and reception by proxy, respectively between a part or all of two or more of said channels, and said communication device, and a logical mirror site is built with each proxy equipment and said communication device to it. Between each proxy equipment and said communication link is performed, and between each proxy equipment and said external terminal, it constitutes so that the gestalt depending on the channel connected to the proxy equipment concerned may be communicated. Said proxy equipment is changing the address of a data transmission place into the transmission place address of self while changing the address of for example, data transmitting origin into the transmitting agency address of self, and it forms the data transceiver gestalt between said external terminals and

said communication devices virtually.

[0012] In the communication system which transmits response data to the addressing to an external terminal concerned through either of two or more channels when the communication system concerning other configurations of this invention receives a data Request to Send from an external terminal Two or more interface devices for corresponding to each of two or more of said channels, and 1 to 1, and enabling the communication link between each channels, While holding the correlation information on a distinction means to distinguish the data characteristic of response data based on said data Request to Send, and each of two or more of said interface devices and a predetermined data characteristic The interface device corresponding to the data characteristic distinguished with said distinction means is specified based on said correlation information, and it is characterized by forming the management tool which leads said response data to said specified interface device.

[0013]

[Embodiment of the Invention] Hereafter, the gestalt of operation of this invention is explained with reference to a drawing.

(The 1st operation gestalt) Drawing 1 is the important section block diagram of the communication system which applied this invention. Among the external terminals (a "client" is called hereafter) 30 and 31 connected to the wide area network 10, for example, the Internet, respectively, a site 300 connects mutually the communication device (a "server" is called hereafter) 310 in which two-way communication is possible, and the proxy server 320 used as an example of proxy (proxy) equipment through an internal network, and is constituted. [0014] The server 310 is equipped with the interface device 311 to which the private address b was assigned, and the interface device 312 to which the global address F was assigned. [0015] It is the substitute communication device equipped with the function to execute transmission and reception of the data between clients 30 and 31 and a server 310 by proxy, a proxy server 320 is changing the address of a data transmission place into the transmission place address of self while it changes the address of data transmitting origin into the transmitting agency address of self, and it forms the data transceiver gestalt between clients 30 and 31 and a server 310 virtually. This proxy server 320 is equipped with the interface device 321 to which the private address c was assigned, and the interface device 322 to which the global address D was assigned.

[0016] In the example of <u>drawing 1</u>, interface devices 311 and 321 perform control to control the internal two-way communication between the server 310-proxy servers 320 through an internal network (local communication link), and for interface devices 312 and 322 each, connect a server 310 and a proxy server 320 to the Internet 10, respectively. In addition, in the communication system of this operation gestalt, the data transmitted and received between a server 310 and each clients 30 and 31 shall be divided under a fixed regulation, and shall be performed in the form of [which the destination label attached] a packet. Hereafter, the data transmitted and received may be called a "packet."

[0017] Actuation of the communication system of this operation gestalt is as follows. Here, the case where the client 30 to which Address A was assigned, and the client 31 to which Address E was assigned publish a data Request to Send towards a server 310 at a coincidence term is explained.

[0018] The data Request to Send from a client 30 shall be published to the proxy server 320 which has an interface device 322. A proxy server 320 will publish a data Request to Send to a server 310 instead of a client 30, if this data Request to Send is received. The publishing agency address of the data Request to Send at this time turns into Address c. That is, the address A of the client 30 which is the issue origin of a data Request to Send from the first is changed into the address c of a proxy server 320, and is sent to a server 310. Although a server 310 transmits a response packet corresponding to this data Request to Send, the transmission place of the response packet at this time serves as a proxy server 320. Once a proxy server 320 changes into Address c the address b by the side of the server 310 which transmitted the response packet, it changes this into the address D of an interface device 322, and transmits the response packet which includes this address D in a header unit to a client 30.

[0019] A server 310 does not need to be conscious of the address A of the issue origin (client 30) of a data Request to Send, and just comes to answer the proxy server 320 by adopting such a gestalt in the data (packet) demanded from the proxy server 320. On the other hand, direct access of the client 31 can be carried out to a server 310 through the interface device 312 with a channel 21 and Address F.

[0020] That is, from the Internet 10 side (a client 30 side, 31 sides), it is visible as if the server with Address D and the server with Address F existed independently, respectively. On the other hand, a local communication link is performed between a server 310 and a proxy server 320. Consequently, the mirror site which has two servers virtually can be built by one set of a server 310, and it becomes possible to distribute the traffic in a server 310 simply and suitable for two channels 20 and 21. Moreover, since it is not necessary to install two sets of servers, a manufacturing cost and maintenance employment expense become cheap.

[0021] (The 2nd operation gestalt) <u>Drawing 2</u> is drawing having shown the example of a configuration of the communication system which formed the above—mentioned proxy server 320 and the proxy servers 420 and 430 of the same function between a server 410 and channels 20 and 21, respectively. One site where the sign 400 arranged this communication system, the interface device (the interface devices 311 and 312 and this function of the 1st operation gestalt) which a server 410 equips with 411 and 412, and 421, 422, 431 and 432 are interface devices (the interface devices 321 and 322 and this function of the 1st operation gestalt) with which each proxy servers 420 and 430 are equipped.

[0022] With the communication system of this operation gestalt as well as the communication system of the 1st operation gestalt, two servers can be virtually built by the single server 410. The other party of data transmission and reception of the server 410 in this case always becomes proxy servers 420 and 430. Consequently, it becomes possible to distribute simply and appropriately the traffic in the server 410 when there is a data Request to Send from two clients 30 and 31.

[0023] (The 3rd operation gestalt) Drawing 3 is the block diagram of the communication system which can distribute traffic, without making a proxy server intervene. Although the channels 20 and 21 which constitute a part use the Internet 10 the same as the thing of the 1st and 2nd operation gestalt and it is not illustrated for convenience, two or more clients shall be connected to the Internet 10. The server 510 used as an example of the communication system in this case equips each of two channels 20 and 21 with the interface devices 511 and 512 corresponding to 1 to 1. The respectively global address is assigned to these interface devices 511 and 512. [0024] First, the structure of the packet (a data Request to Send, response packet) used in this operation gestalt is explained. Data characteristics, such as a transmitting agency, and a communications protocol besides the address of a transmission place, are usually included in the header unit of this kind of packet as additional information. Drawing 4 (a) - (c) is drawing which illustrated the example of a frame structure of a response packet, and the response packet to which Packets A and C are transmitted through an interface device 511 (address 255.20.34.12), and Packet B are packets transmitted through an interface device 512 (address: 255.20.34.13). The transmission place of these packet A-C serves as a client (address: 155.19.25.37, 175.20.35.10, 185.45.67.21) different, respectively, and, as for HTTP and Packet B, also in the classification of a communications protocol, Packets A and C serve as FTP. [0025] The thing [priority] according to a packet is recorded. A "priority" is an index which shows the significance (or quality) of the response for example, according to the contents of service. That is, it is not to concern a server 510 with reception sequence, but transmit a response packet sequentially from what has a high priority to the data Request to Send received within a fixed period. In addition, the same header unit is added also about the data Request to Send, and it is used in the various distinction sections 513-516 mentioned later for the

[0026] The output interface management section 17 performs control for choosing either of the interface devices 511 and 512 at the time of management of routing table 518, and transmission of the response packet according to the data Request to Send from a client, and drawing a response packet at it. Selection of interface devices 511 and 512 is performed based on the

specification of a data characteristic.

newest contents of record of routing table 518 in principle.

[0027] Routing table 518 records the correlation information on each of interface devices 511 and 512, and a predetermined data characteristic in consideration of the engine performance of channels 20 and 21, quality, etc. <u>Drawing 5</u> is drawing having shown the example of the contents of record of this routing table 518, and the data characteristic distinguished in the various distinction sections 513–516 mentioned later besides the transmission place address and the transmitting agency address is recorded. An "output interface device" is an interface device beforehand associated corresponding to each data characteristic.

[0028] The protocol distinction section 513 distinguishes the communications protocol used between clients as a data characteristic. If this communications protocol is FTP and HTTP, it is a mass channel and a data characteristic to choose the channel of small capacity, if it becomes other than this. The packet size distinction section 514 distinguishes the size of a packet size (data size) as a data characteristic. A packet size is a data characteristic to choose a mass channel, so that it is long. The output priority distinction section 515 distinguishes the priority of the response for example, corresponding to a data Request to Send as a data characteristic. This priority determines the sequence of a response according to the importance (or height of quality) of service as mentioned above, and also the thing which has a high priority can use it as a data characteristic to choose a more nearly mass channel. 1 byte of low order of a destination address [in / in the data bit distinction section 516 / the contents of data of the specific location of data, for example, a header unit,] distinguishes odd number or even number as said data characteristic.

[0029] These data characteristics can usually be distinguished in analyzing the contents of the header unit in the case of a data Request to Send. This is because the response packet of the contents according to the data Request to Send will be generated. However, you may make it distinguish the above-mentioned data characteristic from the header unit of a response packet. The distinction result by each distinction sections 513–516 is notified to the output interface management section 517, respectively, and the contents are recorded on routing table 518. In addition, the contents of the above-mentioned data characteristic are instantiation, and are not limited to these things.

[0030] Thus, in the communication system constituted, actuation when there is a data Request to Send from a certain client is as follows. The information on the header unit of a data Request to Send is sent to each distinction sections 513–516, and the data characteristic of the response packet which corresponds, respectively is distinguished. A distinction result is recorded on routing table 518 through the output interface management section 517. If a response packet is generated by the response packet generation means which is not illustrated, the output interface management section 517 will determine to which of interface devices 511 and 512 this response packet is led.

[0031] Although the contents of record of routing table 518 are followed as mentioned above in principle on the occasion of decision, it can also be based only on the distinction result of each distinction sections 513–516. For example, if the number of the tails of the address of a transmission place is odd, in the case of an interface device 511 and even number, an interface device 512 will be determined as an interface device for transmission of a response packet. Or if the protocol of a response packet is HTTP or FTP, an interface device 511 and when other, an interface device 512 will be determined as an interface device for transmission of a response packet. Or when a priority is "quantity" and an interface device 511 and a transmitting priority are "low", an interface device 512 is determined as an interface device for transmission of a response packet. Or the method of determining an interface device by combining suitably the address of the decision and the transmission place according to the merits and demerits of response packet length, a communications protocol, a priority, data size, etc. is also one of the gestalten of operation. According to this communication system, the traffic in a server 510 can be appropriately distributed now.

[0032] the communication link which the response (turn around time) of a server will improve and will be performed through the Internet by building the communication system of the 1st - the 3rd operation gestalt as mentioned above if it sees from a user -- breaking off (quality

fluctuation) — the decreasing effectiveness is acquired, if it sees on the other hand from the side which provides this communication system with service information, the installation number of a server becomes fewer and the effectiveness that the cost which plant—and—equipment investment takes can reduce remarkably will be acquired. In order to build a mirror site logically by one set of a server, it becomes unnecessary moreover, to take the synchronization of an updating stage when it becomes unnecessary to have reproduced the contents in the case of building a mirror site by two sets of servers and there is modification of contents.

[0033] In addition, although the 1st — the 3rd operation gestalt showed the example in case the number of channels is two, this invention is applicable similarly about the case (in this case, the interface device corresponding to 1 to 1 is needed for each channel) of three or more channels.

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DESCRIPTION OF DRAWINGS

[Brief Description of the Drawings]

[Drawing 1] The block diagram of the 1st communication system which applied this invention.

[Drawing 2] The block diagram of the 2nd communication system which applied this invention.

[Drawing 3] The block diagram of the 3rd communication system which applied this invention.

[Drawing 4] (a) - (c) is the explanatory view having shown the example of a configuration of the packet transmitted and received with communication system.

[Drawing 5] The explanatory view having shown the example of a configuration of routing table.

[Drawing 6] The block diagram of the conventional communication system which prepared two servers

[Drawing 7] The block diagram of the conventional communication system equipped with two interface devices.

[Description of Notations]

10 Internet

20 21 Channel

30 31 Client

100 110,200,300,400 Site

320, 420, 430 Proxy server

101, 111, 210, 310, 410, 510 servers

102, 112, 211, 212, 321, 322, 311, 312, 411, 412, 421, 422, 431, 432, 511, and 512 Interface device

513 Protocol Distinction Section

514 Packet Size Distinction Section

515 Priority Distinction Section

516 Data Bit Distinction Section

(19)日本国特許庁 (JP) (12) 公開特許公報 (A)

(11)特許出願公開番号 特開2000-261499 (P2000-261499A)

(43)公開日 平成12年9月22日(2000.9.22)

(51) Int.Cl.7		識別記号	FΙ		テーマコード(参考)
H04L	12/56		H04L 11	1/20 1 0 2 D	5B089
G06F	13/00	3 5 7	G06F 13	3/00 3 5 7 Z	5 K O 3 O
			H04L 11	1/20 1 0 2 E	

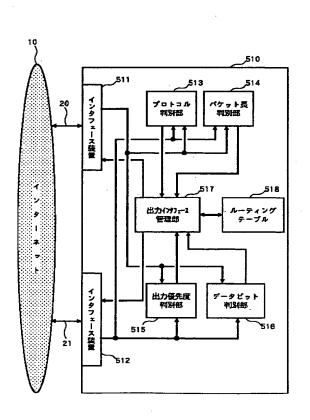
		審查請求	未請求 請求項の数11 OL (全 9 頁)
(21)出願番号	特願平11-59337	(71) 出願人	000102728
			株式会社エヌ・ティ・ティ・データ
(22)出願日	平成11年3月5日(1999.3.5)		東京都江東区豊洲三丁目3番3号
		(72)発明者	藤本 浩
			東京都江東区豊洲三丁目3番3号 株式会
			社工ヌ・ティ・ティ・データ内
		(72)発明者	神崎洋
			東京都江東区豊洲三丁目3番3号 株式会
			社エヌ・ティ・ティ・データ内
•		(74)代理人	100099324
•			弁理士 鈴木 正剛
	•		
•			最終百に続く

(54) 【発明の名称】 トラヒック分散方法及び通信システム

(57)【要約】

【課題】 簡易な手法でトラヒックを分散することがで きる通信システムを提供する。

【解決手段】 複数のインタフェース装置511、51 2と、インターネット10からのデータ送信要求に基づ いてデータ特性を判別する各種データ特性判別部513 ~516、及び、各データ特性判別部513~516で 判別したデータ特性とインタフェース装置511、51 2とを関連付けることにより、どのインタフェース装置 を介して応答データを送信するかを決定する送信インタ フェース管理部517を備えて通信システムを構築す る。



【特許請求の範囲】

【請求項1】 外部端末との間で複数の通信路のいずれかを選択的に使用してデータ送受信を行う装置において実行される方法であって、

前記複数の通信路の一部又は全部と前記装置との間に、 それぞれ前記データ送受信に関する処理を代行するプロ キシ装置を介在させて論理的なミラーサイトを構築し、 個々のプロキシ装置と前記装置との間ではローカルな通 信を行い、各プロキシ装置と前記外部端末との間では当 該プロキシ装置に接続される通信路に依存した形態の通 信を行うことにより、当該装置におけるトラヒックを分 散させることを特徴とする、

トラヒック分散方法。

【請求項2】 外部端末からデータ送信要求を受信したときに当該外部端末宛に複数の通信路のいずれかを介して応答データを送信する装置において実行される方法であって、

前記複数の通信路の各々を所定のデータ特性と関連付けておき、前記外部端末への応答データの送信の際に当該応答データのデータ特性を特定し、この特定したデータ特性に適合する通信路へ前記応答データを導くことにより、前記装置におけるトラヒックを分散させることを特徴とする、トラヒック分散方法。

【請求項3】 前記所定のデータ特性が、前記外部端末との間で使用される通信プロトコル、前記応答データのサイズ、特定位置のデータ内容、応答の優先度、の少なくとも一つを含み、個々の通信路の性能又は品質に応じて決定される情報であることを特徴とする、

請求項2記載のトラヒック分散方法。

【請求項4】 前記受信したデータ送信要求の内容に基づいて、送信すべき前記応答データのデータ特性を特定することを特徴とする、

請求項2又は3記載のトラヒック分散方法。

【請求項5】 前記複数の通信路の一部又は全部と前記 装置との間に、それぞれ前記データ送受信に関する処理 を代行するプロキシ装置を介在させ、少なくとも前記外 部端末への応答データの送信を前記プロキシ装置を通じて行うことを特徴とする、

請求項2又は3記載のトラヒック分散方法。

【請求項6】 外部端末との間で複数の通信路のいずれかを選択的に使用してデータ送受信を行う通信装置を備えた通信システムにおいて、

前記複数の通信路の一部又は全部と前記通信装置との間に、それぞれ前記データ送受信に関する処理を代行する プロキシ装置を配置して各プロキシ装置と前記通信装置 とで論理的なミラーサイトを構築し、

個々のプロキシ装置と前記通信装置との間ではローカルな通信を行い、各プロキシ装置と前記外部端末との間では当該プロキシ装置に接続される通信路に依存した形態の通信を行うように構成したことを特徴とする、通信シ

ステム。

【請求項7】 外部端末からデータ送信要求を受信した ときに当該外部端末宛に複数の通信路のいずれかを介し て応答データを送信する通信システムにおいて、

前記複数の通信路の各々と1対1に対応し個々の通信路 との間の通信を可能にするための複数のインタフェース 装置と、

前記データ送信要求に基づく応答データのデータ特性を 判別する判別手段と、

前記複数のインタフェース装置の各々と所定のデータ特性との関連付け情報を保持するとともに、前記判別手段で判別されたデータ特性に対応するインタフェース装置を前記関連付け情報をもとに特定し、前記応答データを前記特定したインタフェース装置に導く管理手段と、を設けたことを特徴とする通信システム。

【請求項8】 前記判別手段は、前記外部端末との間で使用される通信プロトコルを前記データ特性として判別するものであり、前記管理手段は、前記判別された通信プロトコルに準拠した通信路に接続されるインタフェース装置を特定することを特徴とする、

請求項7記載の通信システム。

【請求項9】 前記判別手段は、前記応答データのサイズの大小を前記データ特性として判別するものであり、前記管理手段は、前記判別されたサイズが大きい順に、大容量の通信路に接続されるインタフェース装置を特定することを特徴とする、

請求項7記載の通信システム。

【請求項10】 前記判別手段は、前記応答データの特定位置のデータ内容を前記データ特性として判別するものであり、前記管理手段は、前記判別されたデータ内容毎に予め割り当てた通信路に接続されるインタフェース装置を特定することを特徴とする、

請求項7記載の通信システム。

【請求項11】 前記判別手段は、前記データ送信要求に対応する応答の優先度を前記データ特性として判別するものであり、前記管理手段は、前記判別された優先度の高い順に、大容量の通信路に接続されるインタフェース装置を特定することを特徴とする、

請求項7記載の通信システム。

【発明の詳細な説明】

[0001]

【発明が属する技術分野】この発明は、例えばインターネットに接続されたサーバにおけるトラヒックの集中を 抑制する技術に関する。

[0002]

【発明の背景】インターネットのWWW(World Wide Web)、FTP(File Transfer Protocol)等により情報提供を行うサービスは、ネットワーク・サービス・プロバイダ(Network Service Provider、以下、「プロバイダ」と略称する)を通じて行われるのが一般的である。

プロバイダが行うサービスは、主として、情報を蓄積しておくサーバにユーザが任意にアクセスすることによって行われる。このため、プロバイダのサーバにユーザからのアクセスが集中すると、そのプロバイダが用意している通信路においてトラヒックの渋滞が発生したり、情報を蓄積しているサーバとプロバイダとの間の通信路にトラヒックの渋滞が発生したりする場合がある。このような渋滞が発生すると、ユーザは、スムーズな情報の享受ができなくなる。

【0003】そこで、トラヒックの渋滞を解消するた め、従来、複数のサーバに、共にユーザがアクセス可能 な同一の情報(「ミラー情報」)を蓄積しておくこと で、トラヒックを分散させることが行われている。ミラ 一情報の提供を行う場所は、一般に、「ミラーサイト」 と呼ばれている。ミラーサイトは、図6に示されるよう に、例えばインターネット10に、複数のサイト、より 単純には二つのサイト100、110が接続された形態 となっている。サイト100にはサーバ101、サイト 110にはサーバ111がそれぞれ配備されている。各 サーバ101、111には、それぞれミラー情報が蓄積 されており、さらに、内部ネットワークと通信路20、 21との接続等を制御するためのインタフェース装置1 02、112が配備されている。このようなミラーサイ トを設けることにより、ユーザからのアクセスが一つの サーバに集中することがなくなり、トラフィックの分散 が可能になる。しかしながら、ミラーサイトを設ける と、独立したサイト(サーバ)が少なくとも2系統、つ まり、サーバが少なくとも2台必要となり、システムの 構築や維持管理に要するコストが大きくなってしまう。 【0004】このような問題の解決手段として、1台の サーバに複数の通信路のいずれかを選択的に接続するこ とができるインタフェース装置、例えばルータを設ける ことで、当該サーバにおけるトラヒックを分散させる形 態が考えられる。つまり、図7に示されるように、サイ ト200に用意されている1台のサーバ210に、自己 とインターネット10とを2つの通信路20、21のい ずれかを介して接続できるようにするためのインタフェ ース装置211、212を配備してトラヒックを分散さ せる。

【0005】しかし、図示のように単にインタフェース装置211、212を複数配備しただけでは、必ずしもトラヒックが通信路20、21にうまく分散されるとは限らない。これは、サーバ210に使用されるオペレーティングシステム(Operating System)の多くが、ユーザから送信されてくるデータを受信するときに、通過したデータの情報を記憶しておかないため、サーバ210がユーザの要求に応じてデータを送信する際に、その適切なインタフェース装置を正しく特定できないためである。

【0006】つまり、ルータを備えるサーバにおいて、

いずれの通信路を用いてデータを送信するかは、サーバが備えるルーティングテーブルの記録内容により決まる。例えば、サーバ210が管理するIP(internet protocol)アドレスが割り付けられた通信装置へのデータ送信であれば、ルーティングテーブルを参照して、その通信装置に接続されている通信路を特定することができる。しかし、不特定のアドレス、例えばグローバルアドレスが割り付けられた通信装置からのデータ送信要求に応答するデータを送信する場合、サイト210は、ルーティングテーブルにその通信装置のIPアドレスを記録しておくことができないため、応答するデータを送信する際に使用する通信路(つまりインタフェース装置)を特定することができない。このため、トラヒックをうまく分散できない。

【0007】そこで本発明は、簡易な手法でトラヒックの分散を可能にするトラフィック制御方法を提供することを課題とする。本発明の他の課題は、トラヒック分散機能を備えた低廉な通信システムを提供することにある。

[8000]

【課題を解決するための手段】上記課題を解決する第1発明のトラヒック制御方法は、外部端末との間で複数の通信路のいずれかを選択的に使用してデータ送受信を行う装置において実行される方法であって、前記複数の通信路の一部又は全部と前記装置との間に、それぞれ前記データ送受信に関する処理を代行するプロキシ装置を前記装置との間ではローカルな通信を行い、各プロキシ装置と前記外部端末との間では当該プロキシ装置に接続される通信路に依存した形態の通信を行うことにより、当該装置におけるトラヒックを分散させることを特徴とする。

【0009】第2発明のトラヒック制御方法は、外部端末からデータ送信要求を受信したときに当該外部端末宛に複数の通信路のいずれかを介して応答データを送信する装置において実行される方法であって、前記複数の通信路の各々を所定のデータ特性と関連付けておき、前記外部端末への応答データの送信の際に当該応答データのデータ特性を特定し、この特定したデータ特性に適合する通信路へ前記応答データを導くことにより、前記装置におけるトラヒックを分散させることを特徴とする。

【0010】ここで、所定のデータ特性とは、例えば、前記外部端末との間で使用される通信プロトコル、前記応答データのサイズ、特定位置のデータ内容、応答の優先度、の少なくとも一つを含み、個々の通信路の性能又は品質に応じて決定される情報である。このデータ特性は、応答データのヘッダ等から特定しても良く、あるいは、前記受信したデータ送信要求の内容に基づいて特定するようにしても良い。また、前記複数の通信路の一部又は全部と前記装置との間に、それぞれ前記データ送受

信に関する処理を代行するプロキシ装置を介在させ、少なくとも前記外部端末への応答データの送信を前記プロキシ装置を通じて行うようにしても良い。

【0011】上記他の課題を解決する本発明の通信シス テムは、外部端末との間で複数の通信路のいずれかを選 択的に使用してデータ送受信を行う通信装置を備えた通 信システムにおいて、前記複数の通信路の一部又は全部 と前記通信装置との間にそれぞれ前記データ送受信に関 する処理を代行するプロキシ装置を配置して各プロキシ 装置と前記通信装置とで論理的なミラーサイトを構築 し、個々のプロキシ装置と前記通信装置との間ではロー カルな通信を行い、各プロキシ装置と前記外部端末との 間では当該プロキシ装置に接続される通信路に依存した 形態の通信を行うように構成したものである。前記プロ キシ装置は、例えば、データ送信元のアドレスを自己の 送信元アドレスに変換するとともにデータ送信先のアド レスを自己の送信先アドレスに変換することで、仮想的 に前記外部端末と前記通信装置との間のデータ送受信形 態を形成するものである。

【0012】本発明の他の構成に係る通信システムは、外部端末からデータ送信要求を受信したときに当該外部端末宛に複数の通信路のいずれかを介して応答データを送信する通信システムにおいて、前記複数の通信路の各々と1対1に対応し個々の通信路との間の通信を可能にするための複数のインタフェース装置と、前記データ送信要求に基づく応答データのデータ特性を判別する判別手段と、前記複数のインタフェース装置の各々と所定のデータ特性との関連付け情報を保持するとともに、前記判別手段で判別されたデータ特性に対応するインタフェース装置を前記関連付け情報をもとに特定し、前記応答データを前記特定したインタフェース装置に導く管理手段とを設けたことを特徴とする。

[0013]

【発明の実施の形態】以下、図面を参照して本発明の実施の形態を説明する。

(第1実施形態)図1は、本発明を適用した通信システムの要部構成図である。サイト300は、それぞれ広域通信網、例えばインターネット10に接続された外部端末(以下、「クライアント」と称する)30、31との間で双方向通信が可能な通信装置(以下、「サーバ」と称する)310と、プロキシ(proxy)装置の一例となるプロキシサーバ320とを内部ネットワークを介して相互に接続して構成されている。

【0014】サーバ310は、プライベートなアドレス bが割り付けられたインタフェース装置311と、グロ ーバルなアドレスFが割り付けられたインタフェース装 置312とを備えている。

【 0 0 1 5 】 プロキシサーバ3 2 0 は、クライアント3 0、3 1 とサーバ3 1 0 との間のデータの送受信を代行する機能を備えた代理通信装置であり、データ送信元の

アドレスを自己の送信元アドレスに変換するとともにデータ送信先のアドレスを自己の送信先アドレスに変換することで、仮想的にクライアント30、31とサーバ310との間のデータ送受信形態を形成するものである。このプロキシサーバ320には、プライベートなアドレスcが割り付けられたインタフェース装置321と、グローバルなアドレスDが割り付けられたインタフェース装置322が備えられている。

【0016】図1の例では、インタフェース装置311、321は、それぞれ内部ネットワークを介したサーバ310ープロキシサーバ320間の内部的な双方向通信(ローカル通信)を制御するものであり、インタフェース装置312、322は、サーバ310及びプロキシサーバ320をそれぞれインターネット10に接続するための制御を行うものである。なお、本実施形態の通信システムにおいて、サーバ310と各クライアント30、31との間で送受信されるデータは、一定の規則で区切られ、宛先ラベルのついたパケットの形式で行われるものとする。以下、送受信されるデータを「パケット」と称する場合がある。

【0017】本実施形態の通信システムの動作は、以下のようになる。ここでは、アドレスAが割り当てられたクライアント30と、アドレスEが割り当てられたクライアント31とが同時期にサーバ310に向けてデータ送信要求を発行した場合について説明する。

【0018】クライアント30からのデータ送信要求 は、インタフェース装置322を有するプロキシサーバ 320に対して発行されるものとする。プロキシサーバ 320は、このデータ送信要求を受け付けると、クライ アント30に代わって、サーバ310ヘデータ送信要求 を発行する。このときのデータ送信要求の発行元アドレ スは、アドレスcとなる。即ち、もともとのデータ送信 要求の発行元であるクライアント30のアドレスAが、 プロキシサーバ320のアドレスcに変換され、サーバ 310へ送られる。サーバ310は、このデータ送信要 求に対応して応答パケットを送信するが、このときの応 答パケットの送信先は、プロキシサーバ320となる。 プロキシサーバ320は、応答パケットを送信したサー パ310側のアドレスbをアドレスcに一旦変換した 後、これをインタフェース装置322のアドレスDに変 換し、このアドレスDをヘッダ部に含む応答パケットを クライアント30宛に送信する。

【0019】このような形態を採用することで、サーバ310は、データ送信要求の発行元(クライアント30)のアドレスAを意識する必要がなく、プロキシサーバ320から要求されたデータ(パケット)をそのプロキシサーバ320宛に返信すればよくなる。一方、クライアント31は、通信路21及びアドレスFを持つインタフェース装置312を介してサーバ310に直接アクセスすることができる。

【0020】つまり、インターネット10側(クライアント30、31側)からは、アドレスDを持つサーバと、アドレスFを持つサーバとがそれぞれ独立して存在しているかのように見える。一方、サーバ310とプロキシサーバ320との間ではローカル通信が行われる。この結果、1台のサーバ310で、仮想的に2つのサーバをもつミラーサイトを構築することができ、サーバ310におけるトラヒックを2つの通信路20、21に簡易且つ適切に分散することが可能になる。また、2台のサーバを設置する必要がないので、製造コストや維持運用費が低廉になる。

【0021】(第2実施形態)図2は、サーバ410と通信路20、21との間に、それぞれ上述のプロキシサーバ320と同一機能のプロキシサーバ420、430を設けた通信システムの構成例を示した図である。符号400は、この通信システムを配した一つのサイト、411、412はサーバ410が備えるインタフェース装置(第1実施形態のインタフェース装置311、312と同機能)、421、422、431、432は、各プロキシサーバ420、430が備えるインタフェース装置(第1実施形態のインタフェース装置321、322と同機能)である。

【0022】この実施形態の通信システムでも、第1実施形態の通信システムと同様にして、単一のサーバ41 0により2つのサーバを仮想的に構築することができる。この場合のサーバ410のデータ送受信の相手方は、常にプロキシサーバ420、430となる。その結果、二つのクライアント30、31からデータ送信要求があったときのサーバ410におけるトラヒックを簡易且つ適切に分散することが可能になる。

【0023】(第3実施形態)図3は、プロキシサーバを介在させずにトラヒックを分散することができる通信システムのブロック図である。便宜上、インターネット10を一部を構成する通信路20、21は第1及び第2実施形態のものと同じとし、図示しないが、インターネット10には、複数のクライアントが接続されているものとする。この場合の通信システムの一例となるサーバ510は、2つの通信路20、21の各々に1対1に対応するインタフェース装置511、512を備えている。これらのインタフェース装置511、512には、それぞれグローバルなアドレスが割り付けられている。

【0024】まず、この実施形態において使用されるパケット(データ送信要求、応答パケット)の構造について説明する。この種のパケットのヘッダ部には、通常、送信元や送信先のアドレスのほか、通信プロトコル等のデータ特性が付加情報として含まれる。例えば図4

(a)~(c)は、応答パケットのフレーム構成例を例示した図であり、パケットA、Cはインタフェース装置511(アドレス255.20.34.12)を通じて送信される応答パケット、パケットBはインタフェース装置512

(アドレス: 255. 20. 34. 13) を通じて送信されるパケットである。これらのパケットA~Cの送信先は、それぞれ異なるクライアント(アドレス: 155. 19. 25. 37、175. 20. 35. 10、185. 45. 67. 21) となっており、通信プロトコルの種別も、パケットA、CはHTTP、パケットBはFTPとなっている。

【0025】優先度もパケットに応じたものが記録されるようになっている。「優先度」とは、例えばサービス内容に応じた応答の重要度(あるいは品質)を示す指標である。つまり、サーバ510が一定期間内に受け付けたデータ送信要求に対し、受付順番に関わらず、優先度の高いものから順に応答パケットを送信できるようにするためのものである。なお、データ送信要求についても同様のヘッダ部が付加されており、後述する各種判別部513~516においてデータ特性の特定のために使用されるようになっている。

【0026】出力インタフェース管理部17は、ルーティングテーブル518の管理と、クライアントからのデータ送信要求に応じた応答パケットの送信時にインタフェース装置511、512のいずれかを選択して応答パケットを導くための制御とを行うものである。インタフェース装置511、512の選択は、原則として、ルーティングテーブル518の最新の記録内容に基づいて行われる。

【0027】ルーティングテーブル518は、通信路20、21の性能、品質等を考慮して、インタフェース装置511、512の各々と所定のデータ特性との関連付け情報を記録するものである。図5は、このルーティングテーブル518の記録内容例を示した図であり、送信先アドレス、送信元アドレスのほか、後述する各種判別部513~516で判別されたデータ特性が記録されるようになっている。「出力インタフェース装置」は、各データ特性に対応して予め関連付けられたインタフェース装置である。

【0028】プロトコル判別部513は、クライアント との間で使用される通信プロトコルをデータ特性として 判別するものである。この通信プロトコルは、例えばF TPやHTTPなら大容量の通信路、それ以外ならば小 容量の通信路が選択されるようにするためのデータ特性 である。パケット長判別部514は、例えば、パケット 長(データサイズ)の大小をデータ特性として判別する ものである。パケット長は、それが長いほど大容量の通 信路が選択されるようにするためのデータ特性である。 出力優先度判別部515は、例えば、データ送信要求に 対応する応答の優先度をデータ特性として判別するもの である。この優先度は、前述のようにサービスの重要性 (または品質の高さ)に応じて応答の順序を決定するほ か、優先度が高いものほど、より大容量の通信路が選択 されるようにするためのデータ特性として使用すること ができる。データビット判別部516は、データの特定 位置のデータ内容、例えばヘッダ部における宛先アドレスの下位1パイトが奇数か偶数かを前記データ特性として判別するものである。

【0029】これらのデータ特性は、通常、データ送信要求の際に、そのヘッダ部の内容を解析することで判別が可能である。これは、そのデータ送信要求に応じた内容の応答パケットが生成されることになるからである。但し、応答パケットのヘッダ部から上記データ特性を判別するようにしても良い。各判別部513~516による判別結果は、それぞれ出力インタフェース管理部517に通知され、その内容がルーティングテーブル518に記録されるようになっている。なお、上記データ特性の内容は例示であり、これらのものに限定されるものではない。

【0030】このように構成される通信システムにおいて、あるクライアントからデータ送信要求があったときの動作は、以下のようになる。データ送信要求のヘッダ部の情報が各判別部513~516に送られ、それぞれ対応する応答パケットのデータ特性が判別される。判別結果は、出カインタフェース管理部517を通じてルーティングテーブル518に記録される。図示しない応答パケット生成手段で応答パケットが生成されると、出カインタフェース管理部517は、この応答パケットをインタフェース装置511、512のどちらに導くかを決定する。

【0031】決定に際しては、原則として上述のように ルーティングテーブル518の記録内容に従うが、各判 別部513~516の判別結果のみに基づくこともでき る。例えば、送信先のアドレスの末尾が奇数であればイ ンタフェース装置511、偶数の場合はインタフェース 装置512を応答パケットの送信用のインタフェース装 置として決定する。あるいは、応答パケットのプロトコ ルがHTTP又はFTPであればインタフェース装置5 11、それ以外の場合はインタフェース装置512を応 答パケットの送信用のインタフェース装置として決定す る。あるいは、優先度が"高"の場合はインタフェース 装置511、送信優先度が"低"の場合はインタフェー ス装置512を応答パケットの送信用のインタフェース 装置として決定する。あるいは、応答パケット長の長短 に応じた決定、送信先のアドレス、通信プロトコル、優 先度、データサイズ等を適宜組み合わせることで、イン タフェース装置を決定する方法も実施の形態の一つであ る。この通信システムによれば、サーバ510における トラヒックを適切に分散できるようになる。

【0032】以上のように第1~第3実施形態の通信システムを構築することで、ユーザからみれば、サーバのレスポンス(ターンアラウンドタイム)が向上し、インターネットを通じて行う通信の途切れ(品質変動)が少

なくなる効果が得られ、一方、この通信システムにサービス情報を提供する側からみれば、サーバの設置台数が減り、設備投資に要するコストを著しく低減させることができるという効果が得られる。また、1台のサーバで論理的にミラーサイトを構築するため、2台のサーバでミラーサイトを構築する場合のコンテンツの複製を行う必要なくなり、また、コンテンツの変更があったときの更新時期の同期をとる必要もなくなる。

【0033】なお、第1~第3実施形態では、通信路が2つの場合の例を示したが、3以上の通信路の場合(この場合は各々の通信路に1対1に対応したインタフェース装置が必要になる)についても、同様に、本発明を適用できるものである。

[0034]

【発明の効果】以上の説明から明らかなように、本発明によれば、簡易な手法でトラヒックの分散を図ることができるトラヒック分散方法を提供できるようになる。また、トラヒック分散機能を備えた、低廉な通信システムを提供することができる。

【図面の簡単な説明】

【図1】本発明を適用した第1の通信システムのブロック図。

【図2】本発明を適用した第2の通信システムのブロック図。

【図3】本発明を適用した第3の通信システムのブロック図。

【図4】(a)~(c)は、通信システムで送受信されるパケットの構成例を示した説明図。

【図5】ルーティングテーブルの構成例を示した説明 図。

【図6】2つのサーバを設けた従来の通信システムの構成図。

【図7】2つのインタフェース装置を備えた従来の通信 システムの構成図。

【符号の説明】

10インターネット20、21通信路

30、31 クライアント

100、110、200、300、400 サイト

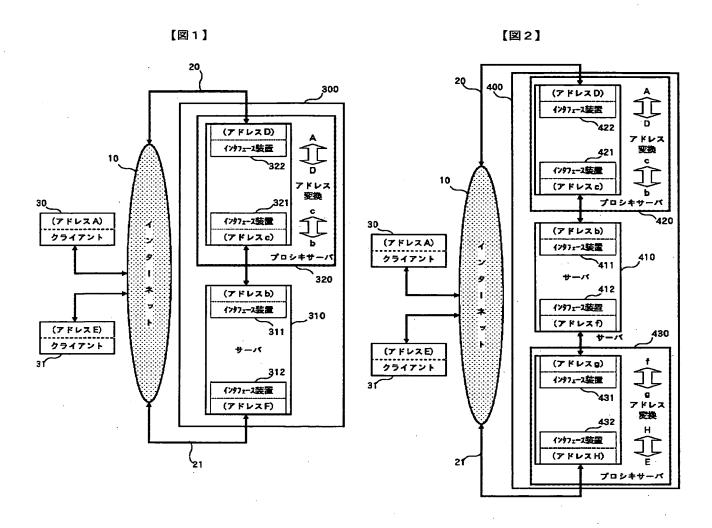
101、111、210、310、410、510サーバ

320、420、430 プロキシサーバ

102、112、211、212、321、322、311、312、411、412、421、422、43 1、432、511、512 インタフェース装置

5 1 3プロトコル判別部5 1 4パケット長判別部5 1 5優先度判別部

516 データビット判別部



【図4】

(a) <パケットA >

		TA/		К		
送信元アドレス (255.20.34.12)	送信先アドレス (155.19.25.37)	プロトコル (HTTP)	優先度 (高)	**	データ	
(b)						
くパケットB>						

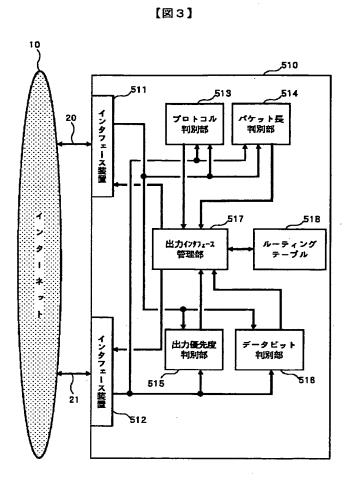
	sc				
送信元アドレス (255.20.34.13)	送信先アドレス (175.20.35.10)	プロトコル (FTP)	優先度 (低)		データ
				- 2)	

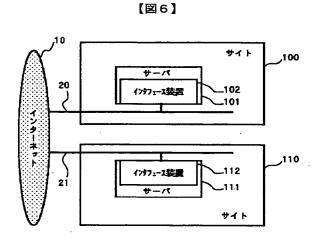
(c)

_	<パケ	» FC>			
送信元アドレス (255.20.34.12)	送信先アドレス (185.45.67.21)	プロトコル (HTTP)	優先度 (高)	sc.	データ

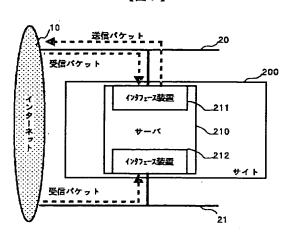
【図5】

送信元アドレス	送信先アドレス	プロトコル	優先度	出力(ンナフュース 装置
255.20.34.12	155.19.25.37	нттр	高	511
255.20.34.13	175.20.35.10	FTP	任	512
255.20.34.12	185.45.67.21	нттр	超	511





【図7】



フロントページの続き

(72)発明者 宮本 久仁男 東京都江東区豊洲三丁目3番3号 株式会 社エヌ・ティ・ティ・データ内

F ターム(参考) 5B089 GA19 GA21 HA10 KA07 KC39 KG05 KG08 KH03 MA03 5K030 GA03 GA13 HA08 HC01 JT02 KA01 KA07 KA13 LB05 LD11 LE01 LE05